# Back to the Futures:

#### Incremental Parallelization of Existing Sequential Runtimes

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- Incremental
- Seq. performance intuition carries over
- Low development investment
  Good scaling (negligible sequential overhead)



### An Observation

- Runtime "fast-path" operations generally have few side effects
- Thus, safe for parallelism

Partition operations into 3 categories:

- Safe (run in parallel)
- Unsafe (runtime side effects)
- A few others (a priori unsafe, but important)

Safety may be dependent on arguments

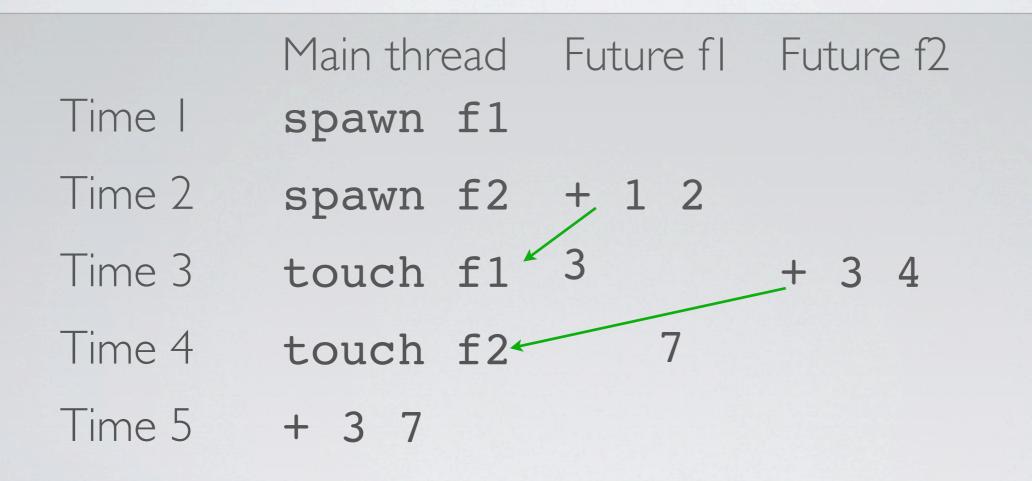
- One runtime thread where everything is safe
- Barricades active on all other threads:
  - Detect and intercept unsafe ops
  - Halt a thread until unsafe op can be completed by runtime thread
- Add primitives allowing programmer to explicitly donate the runtime thread's time to the barricaded thread, allowing it to pass through and continue

### Racket Language Extension

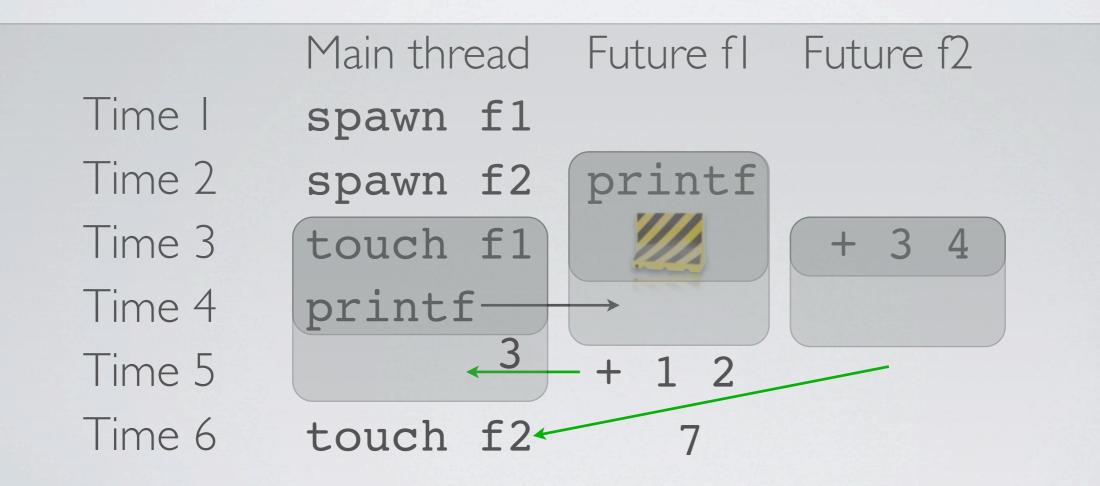
#### future : $(-> \alpha) -> \alpha$ future

#### touch : $\alpha$ future -> $\alpha$

(let ([f1 (future  $(\lambda () (+ 1 2)))]$ [f2 (future  $(\lambda () (+ 3 4)))])$ (+ (touch f1) (touch f2)))



(let ([f1 (future ( $\lambda$  () (printf "Hello!\n") (+ 1 2)))] [f2 (future ( $\lambda$  () (+ 3 4)))]) (+ (touch f1) (touch f2)))



### Racket Implementation

#### Racket runtime:

- Substrate for the Racket language
- I00,000+ lines of C code
- Simple, eager JIT compiler
- Global data includes:
  - Execution state (exception handlers)
  - Symbol table
  - Macro expansion caches
  - GC metadata

# Racket Operations

Safe	Unsafe	Other
+	+	allocation
/	/	JIT compilation
unsafe-fl+	hash-set!	
unsafe-fl/	printf	
unsafe-vector-ref	vector-ref	
unsafe-vector-set!	printf	
	call/cc	
	write	
	read	
	open-input-file	
	error	

#### Barricades in Racket

- All code JIT compiled (if possible)
- Fast-path ops inlined
- Slow-path ops C functions

## "Other" Operations

- We leverage Racket's user-level thread infrastructure for:
  - Allocation
  - JIT compilation
- Racket threads: preemptive to programmers, cooperative to runtime
- Cooperation points allow for polling

### Garbage Collection

 GC = special form of synchronized operation (stop the world)

Cooperation points become barriers

☑ Incremental

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## Development Person-Hours (Racket)

 Performed by non-expert (no prior knowledge) and runtime developer

Expert	Non-Expert (me)	Total
41	536	577

### Parrot Implementation

#### Parrot runtime:

- Register-based virtual machine
- Pluggable runloop allows switching between interpreters
- Dynamic (virtual functions)
- Each bytecode is checked prior to execution for safety
  - Includes argument checking

## Development Person-Hours (Racket)

Performed by expert (active runtime implementation contributor)

Expert	Non-Expert	Total
52	-	52

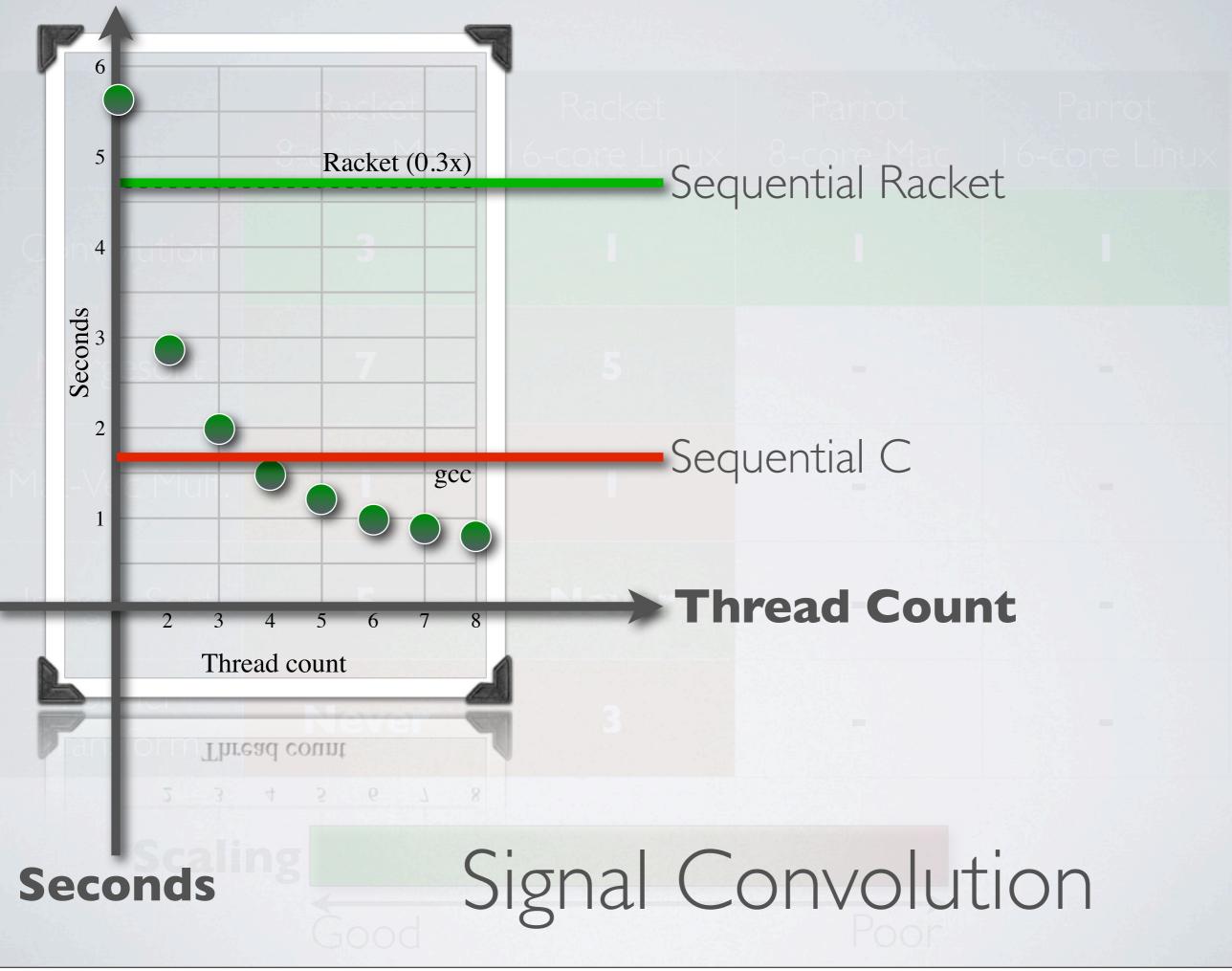
# Performance Evaluation

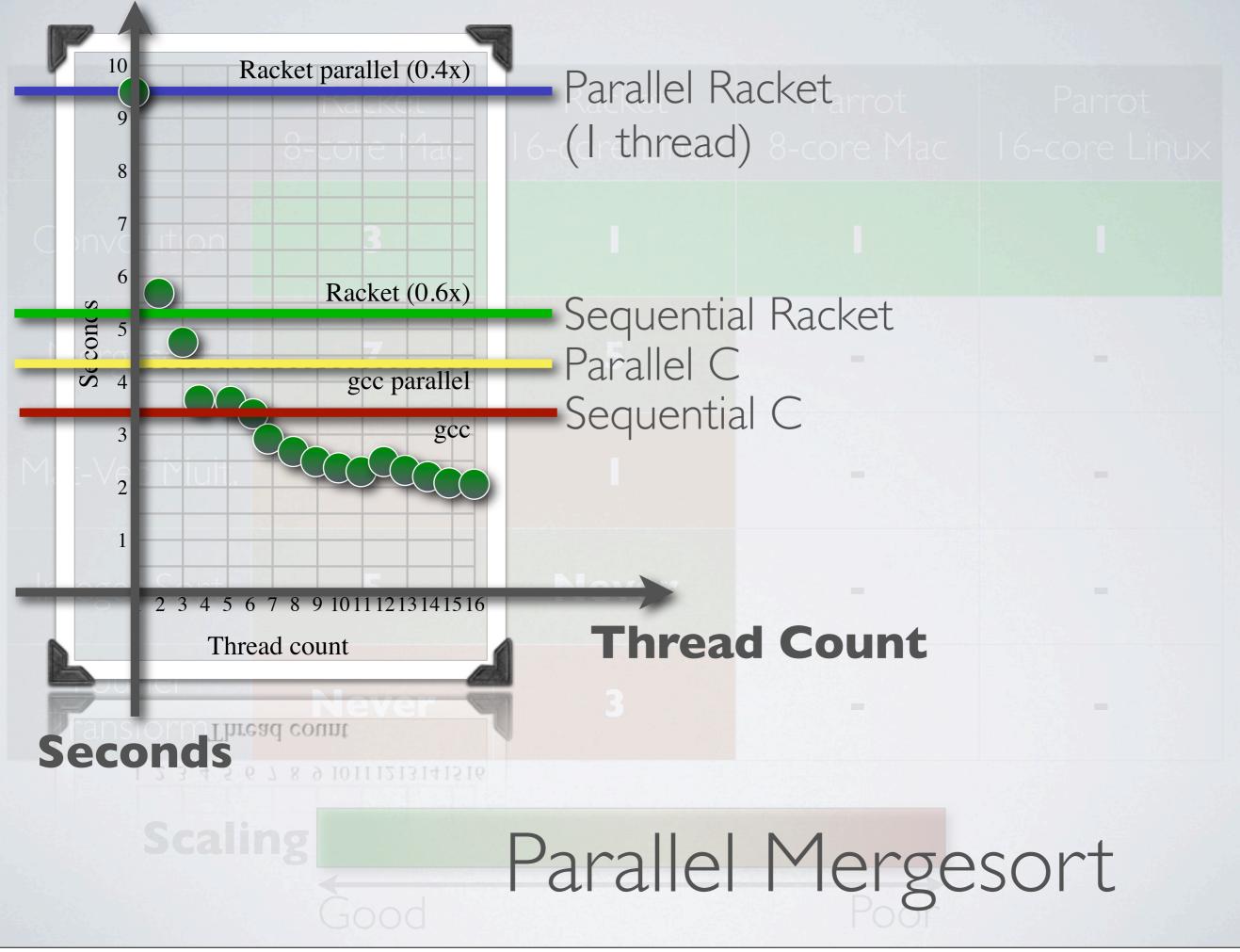
- 3 microbenchmarks
  - Signal convolution
  - Mergesort
  - Sparse matrix-vector multiplication
- 2 NAS Parallel Benchmarks kernels
  - Integer Sort
  - Fourier Transform
- 2 test machines:
  - 8-core workstation (Mac OS X)
  - I6-core mid-range server (Linux)

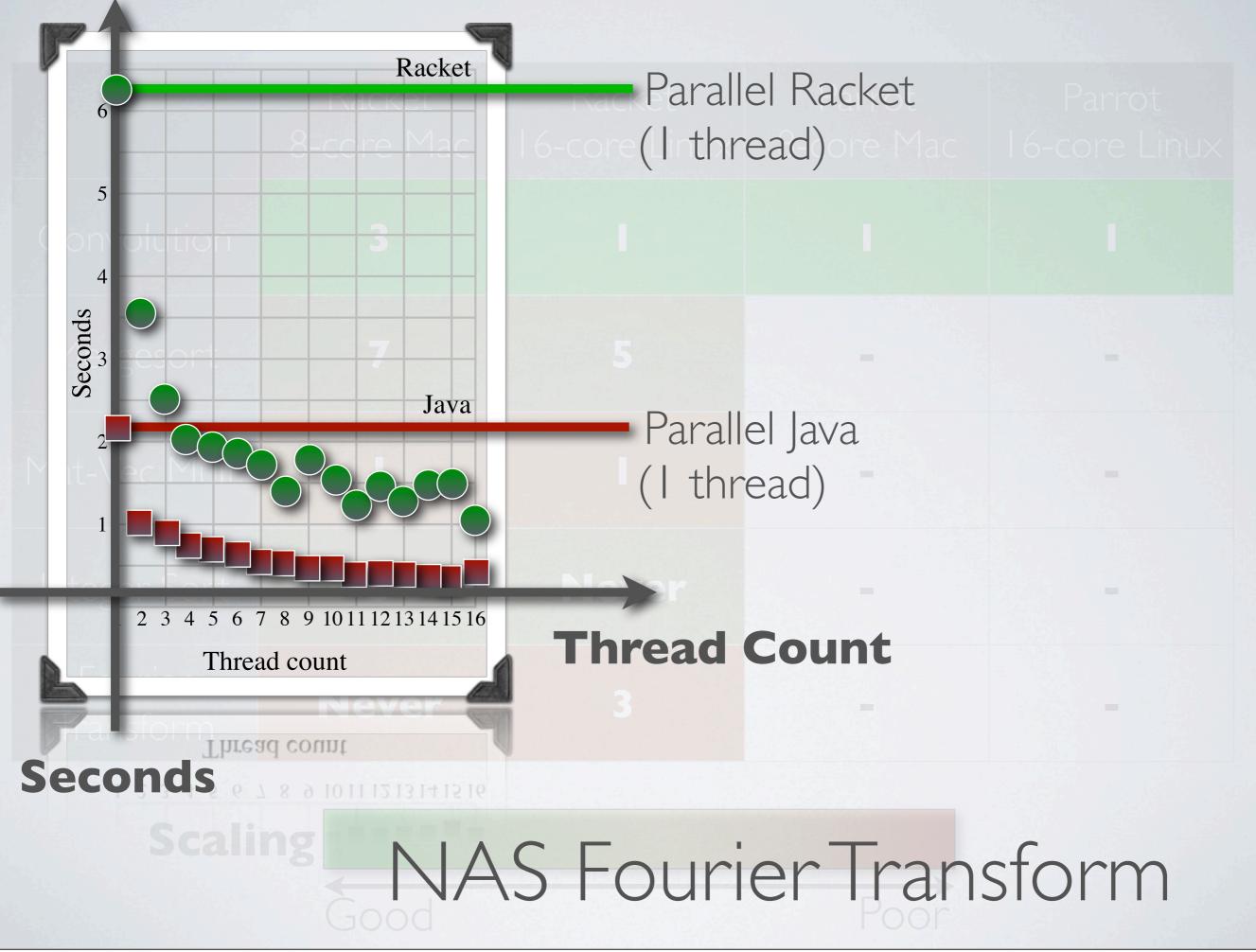
	Racket 8-core Mac	Racket 16-core Linux	Parrot 8-core Mac	Parrot 16-core Linux
Convolution	4	2	2	2
Mergesort	8	6	-	-
Mat-Vec Mult.	2	2	-	-
Integer Sort	6	Never	-	-
Fourier Transform	Never	4	-	-

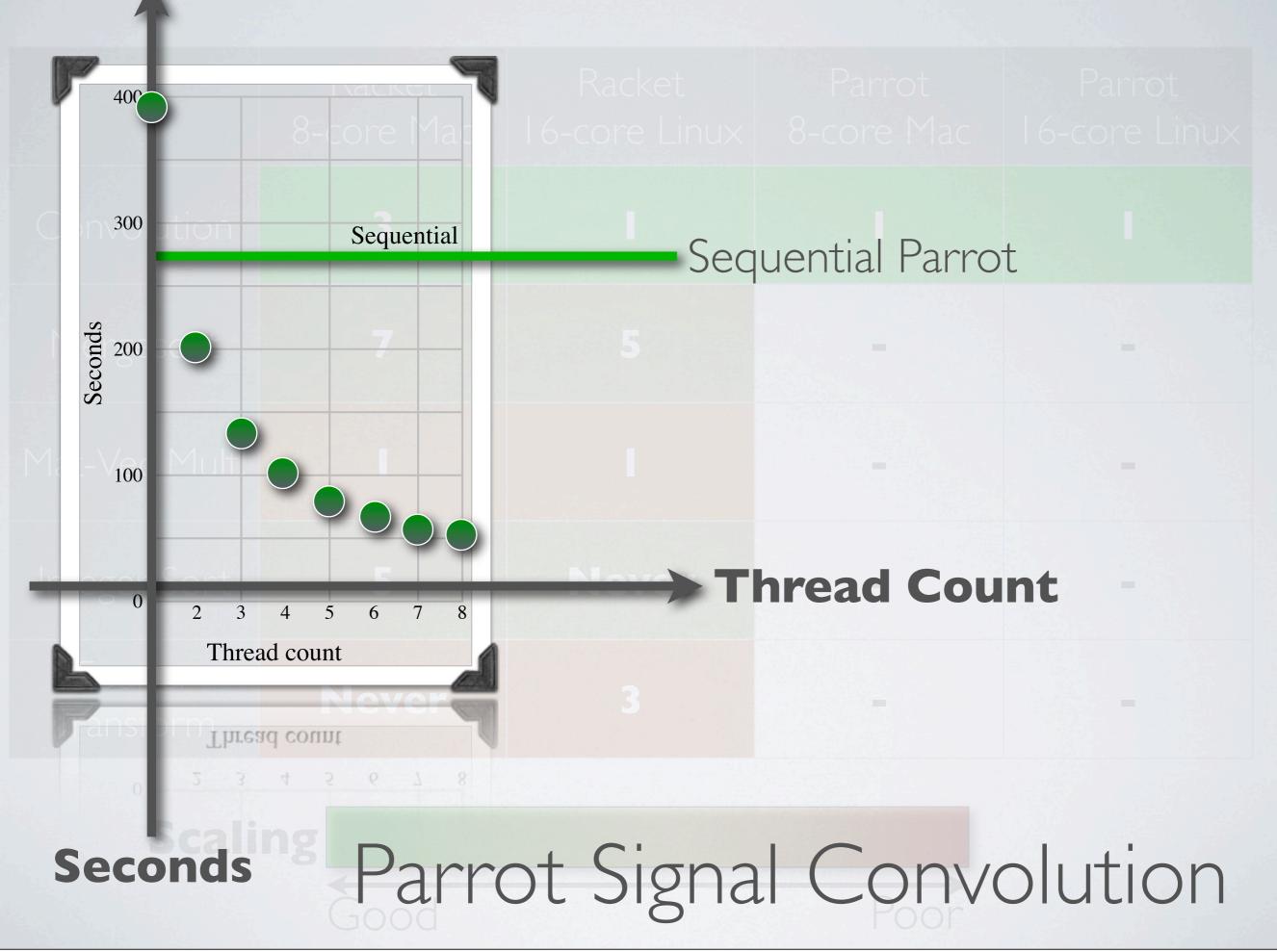


**Values** = # of threads to beat sequential impl.









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#### Thanks!

- Try parallel Racket today: <u>http://racket-lang.org/download/</u>
   Try slow path barricading in your runtime system;
- the main system developer should be able to add it within a few weeks of work