



Opportunities for Leveraging OS Virtualization in High-End Supercomputing

MASVDC'10:

**Workshop on Micro Architectural Support for
Virtualization, Data Center Computing, and Clouds**

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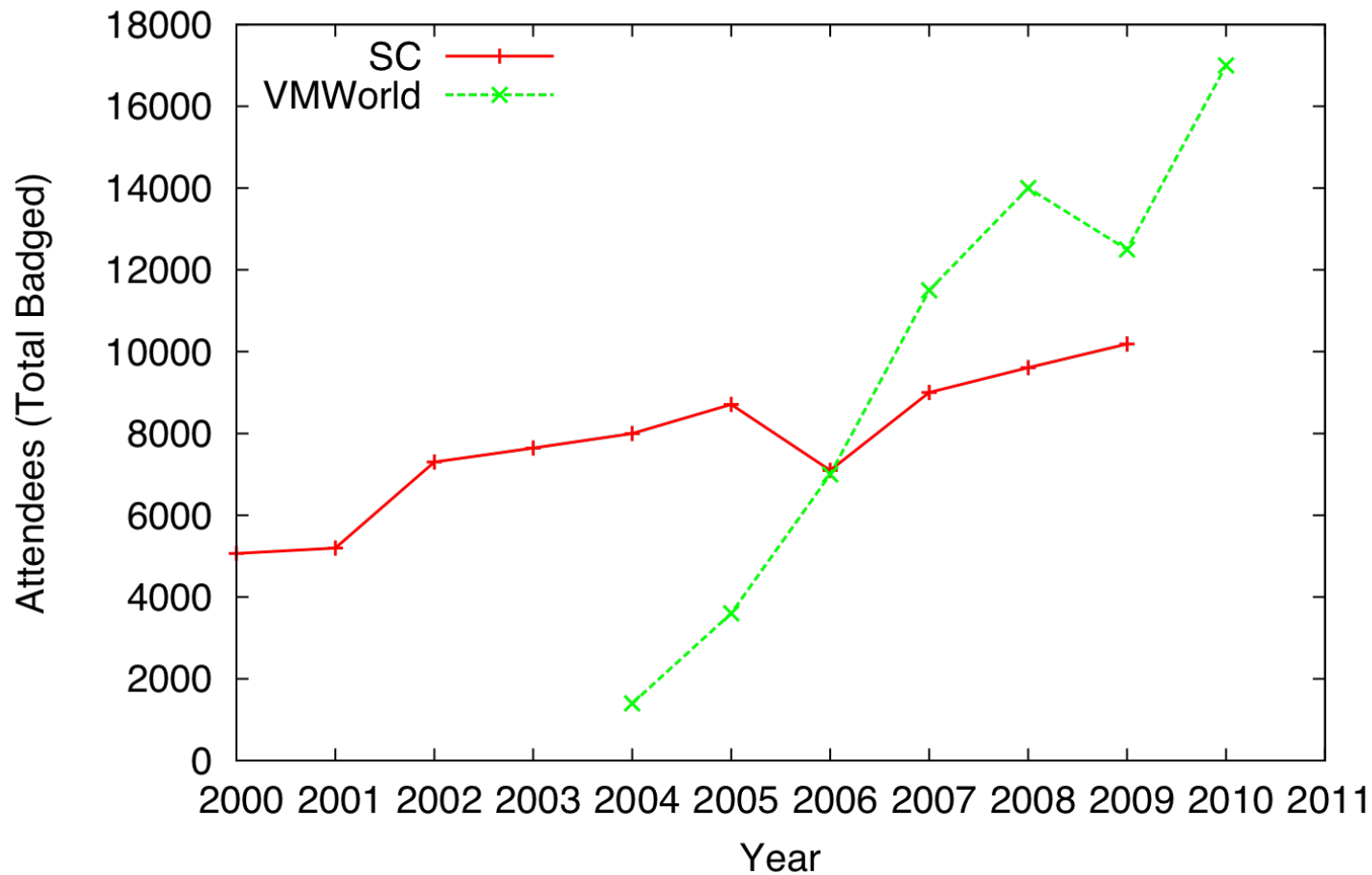
Outline

- **Introduction**
- **Previous Work**
- **High-End HPC Virtualization Use Cases**
- **Results**
- **Conclusion**



Apples and Oranges, But... No Doubt Mainstream Virtualization Seeing Explosive Growth

SC vs. VMWorld Conference Attendance



Sources: SC web sites, news articles, blog posts



Virtualization in HPC?

- “Every problem in computer science can be solved with another level of abstraction” ;-)
- “No virtualization in HPC”
 - Well, we (usually) have virtual memory
 - Virtualization is potentially disruptive
 - Clayton M. Christensen's keynote at SC'10
 - Won't/Can't attack established HPC initially, may sneak up over time

Vendors steadily decreasing virtualization overhead and adding capabilities



Virtualization in High-End HPC?

- **Compelling use cases not necessarily dependent on achieving absolute highest performance**
 - Increase flexibility, app-specific OS/runtime
 - Enable new capabilities not present today
 - Modest overheads tolerable
- **Well known techniques such as VMM-bypass and large paging mitigate overheads**

Our results show virtualization overhead is low, typically less than 5%



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Previous Work: Motivation and I/O Optimization

- **Motivation for migrating HPC workloads to VMs**

(ICS'06: Huang, Liu, Abali, Panda)

- **Ease of management (live migration, checkpoint)**
- **Ability to run custom tailored OS (LWK)**
- **Exposing privileged ops to user (kernel modules)**

- **High-performance I/O**

– **VMM-bypass** *(USENIX'06: Liu, Huang, Abali, Panda)*

– **Migrating VMM-bypass VMs** *(VEE'07: Huang, Liu, Koop, Abali, Panda)*

– **PGAS applications in Xen VMs**

(Cluster'07: Scarpazza, Mullaney, Villa, Petrini, Tipparaju, Brown, Nieplocha)



Previous Work: Resiliency and Overhead Reduction

- **Proactive VM migration to improve resiliency**
(ICS'07: Nagarajan, Mueller, Engelmann, Scott)
(FGCS-Mar10: Scott, Vallee, Naughton, Tikotekar, Engelmann, Ong)
 - Migrate away from nodes with observed deteriorating health
 - Reactive checkpoint frequency can be reduced if MTTI improved

- **Nested paging to reduce VM exits**
 - AMD nested paging, Intel EPT
 - 2-D nested page table caching scheme
(ASPLOS'08: Bhargava, Serebrin, Spadini, Manne)
 - NPT structure does not have to match native
(CAL-Jan10: Hoang, Bae, Lange, Zhang, Dinda, Joseph)



Previous Work: Cloud and VM Scalability

- **Using public clouds for HPC**
 - **Migrating workloads and performance measurements**
(SC'08: Deelman, Singh, Livny, Berriman, Good)
(GC'09: Hill, Humphrey)
 - **Amazon's EC2 HPC instances with 10GigE + GPUs**

- **Scalability of MPI apps in VM on Cray XT**
(IPDPS'10: Lange, Pedretti, Hudson, Dinda, Cui, Xia, Bridges, Gocke, Jaconette, Levenhagen, Brightwell)
 - **Micro-benchmarks and real applications**
 - **Up to ~6K nodes, more on way**



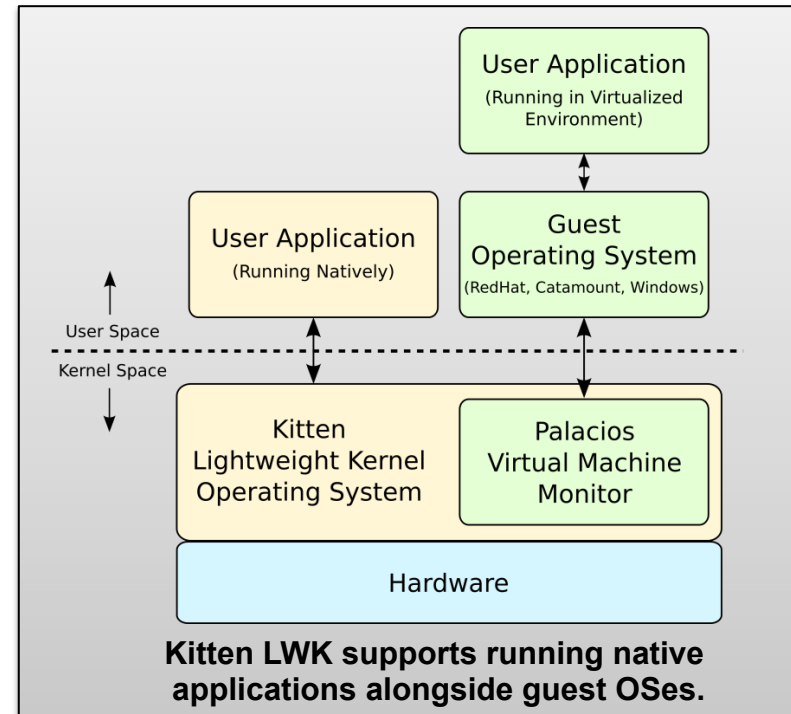
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Enhancing Lightweight OS Flexibility

- Original motivation
- LWK provides high perf. native environment
- VMM allows full-featured guest OS (e.g., Red Hat Linux) to be loaded on-demand
 - Perl, python, matlab, ...
 - COTS databases, simulators, ...
 - You name it
- Approach also applies to lightweight Linux distributions like CLE (Cray Linux Env.)



Kitten available from: <http://code.google.com/p/kitten/>
Palacios available from: <http://v3vee.org/>



Tool for Exascale OS Research

- **Obtaining dedicated time on supercomputer to test prototype OS is HARD**
- **VM capability would partially mitigate**
 - **Test prototype “X-stack” at scale, expose effects that only occur at scale**
 - **Rapid turnaround for debug iterations**
 - **VM is convenient instrumentation layer**
- **Support HW/SW co-design efforts**
 - **Prototype new HW/SW interfaces and capabilities**
 - **Tie to architectural simulator**



Enable New Capabilities

- **Internet-scale simulation**
 - Run commodity OSes and software
 - Multiple virtual nodes per physical node
- **Migration based on VMM-level runtime monitoring**
 - Better map application onto network topology
 - Migrate memory pages among NUMA nodes
 - Make up for all VMM overhead and more (?)
- **Provide backwards compatibility**
 - Support legacy software on future exascale systems
 - Provide incremental path to native environment



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Test Platform

Processor	Intel X5570 2.93 GHz quad-core 2 sockets, 8 cores total 2 NUMA nodes Theoretical Peak: 94 GFLOPS
Memory	24 GB DDR3-1333 Three 4 GB DIMMs per socket Theoretical Peak: 64 GB/s
BIOS Configuration	Hyper-Threading Disabled Turbo-Boost Disabled Maximum Performance
Software	Linux 2.6.36.7 with KVM Guest image identical to host kvm-clock para-virtualized clock, plus ntp daemon NUMA topology exposed to guest libhugetlbfs for large paging



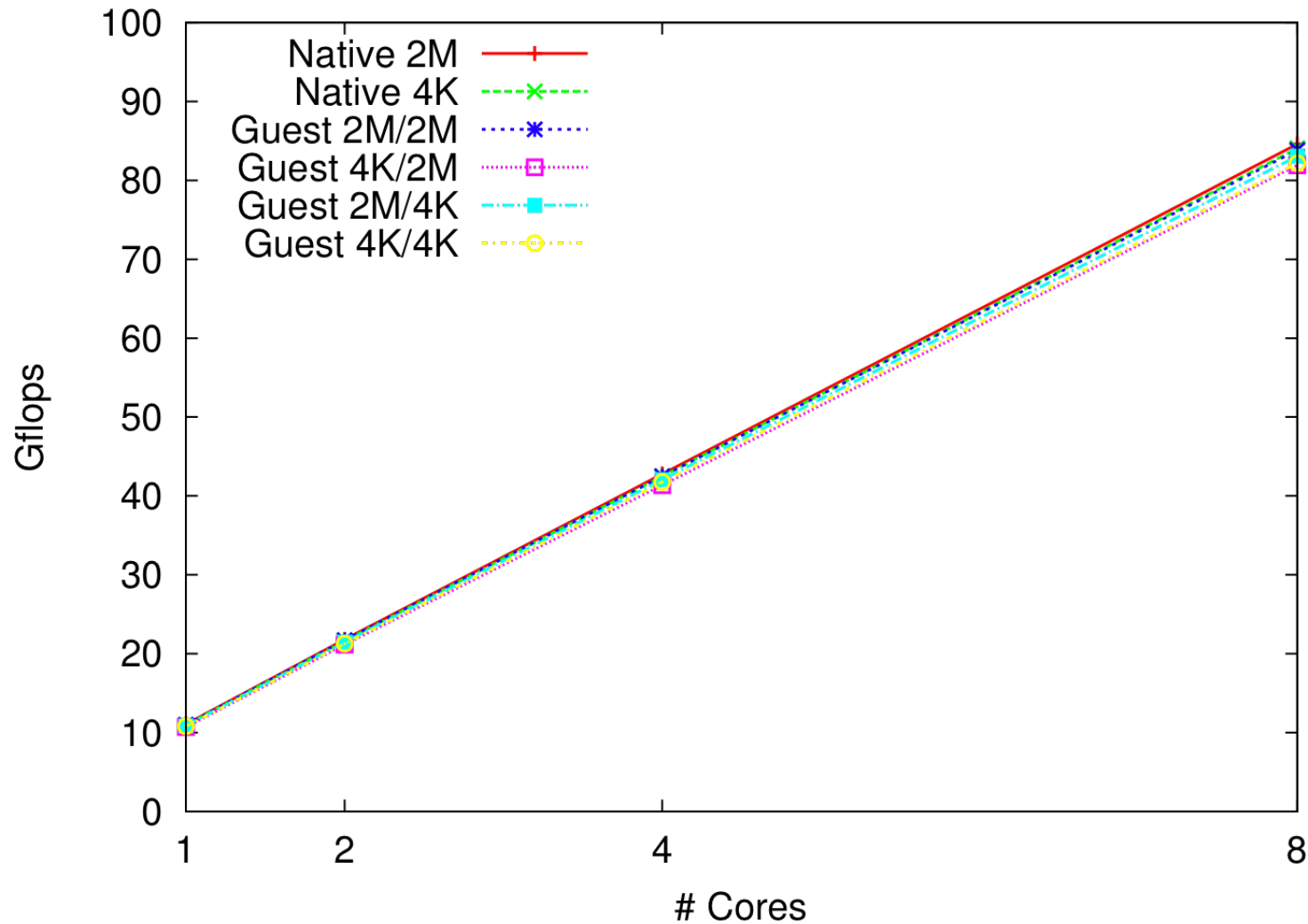
Benchmarks

- **Compute overhead**
 - Linpack (HPCC HPL)
- **Memory overhead**
 - OpenMP STREAM
 - GUPs (HPCC MPIRandomAccess)
- **MPI**
 - PingPong (IMB PingPong)
Intra-node only, via shared mem (MPICH2 Nemesis)



HPL Linpack

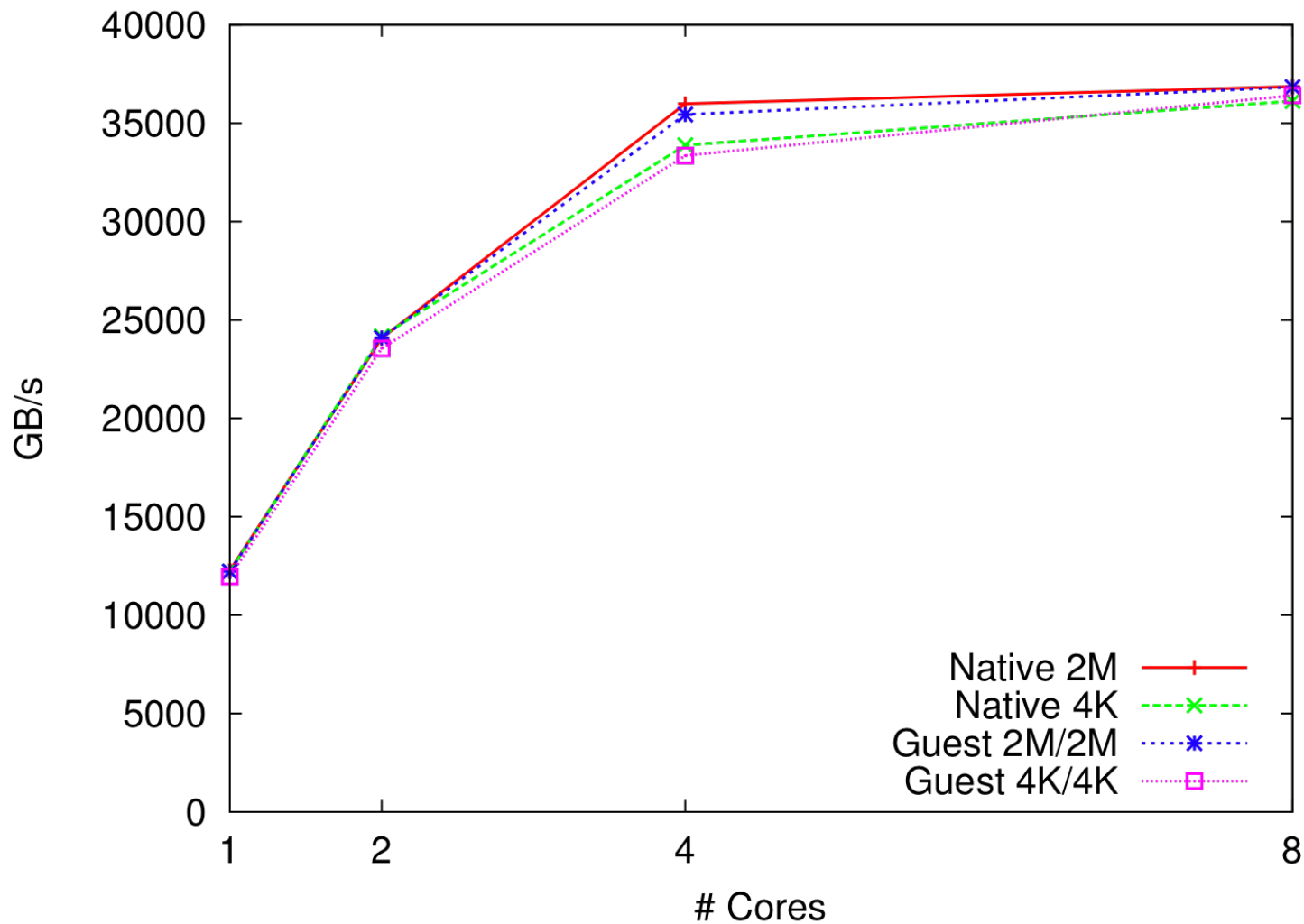
No Compute Virtualization Overhead





OpenMP STREAM

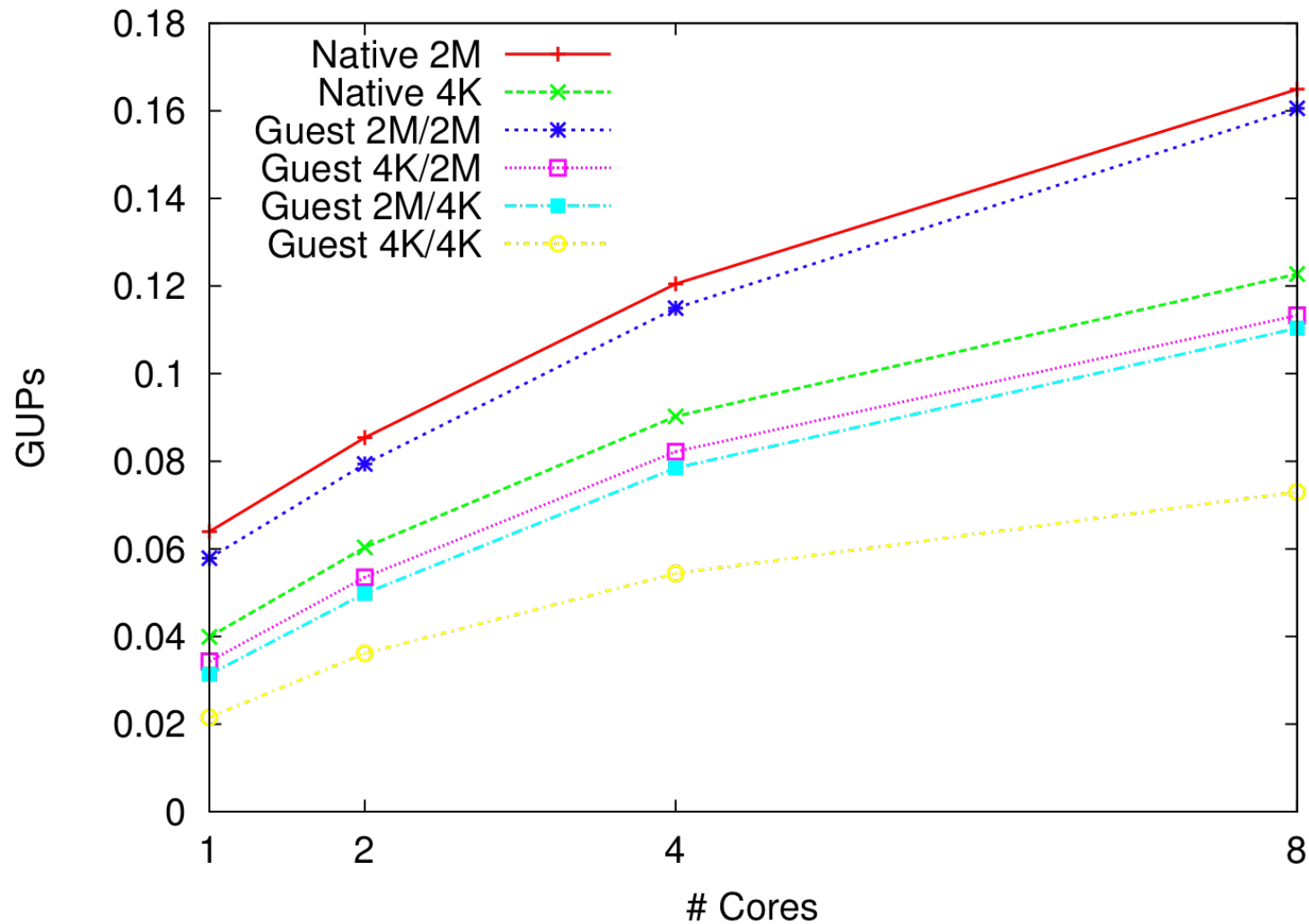
Little Memory BW Virtualization Overhead





MPI Random Access

2.5% to 40% Overhead Depending on Config



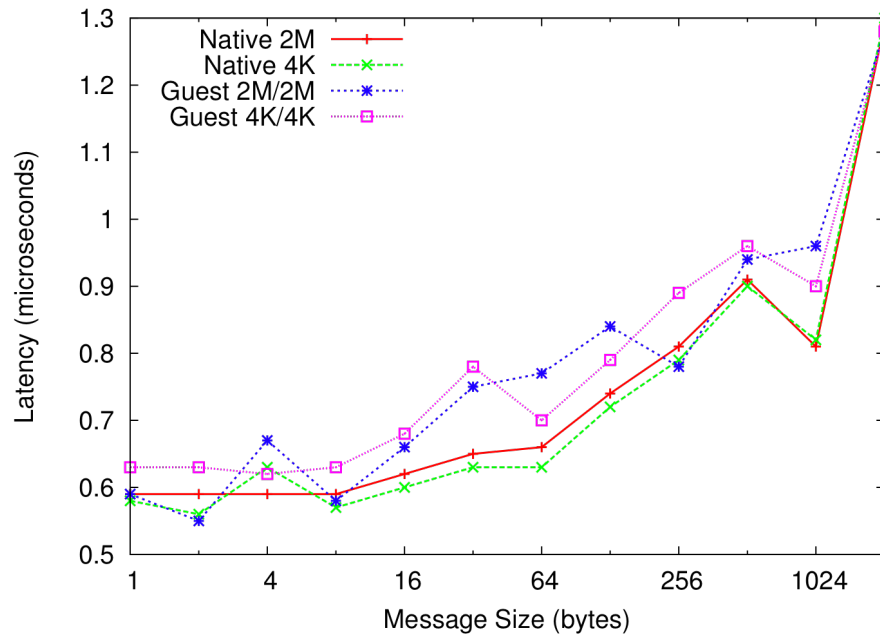


MPI PingPong

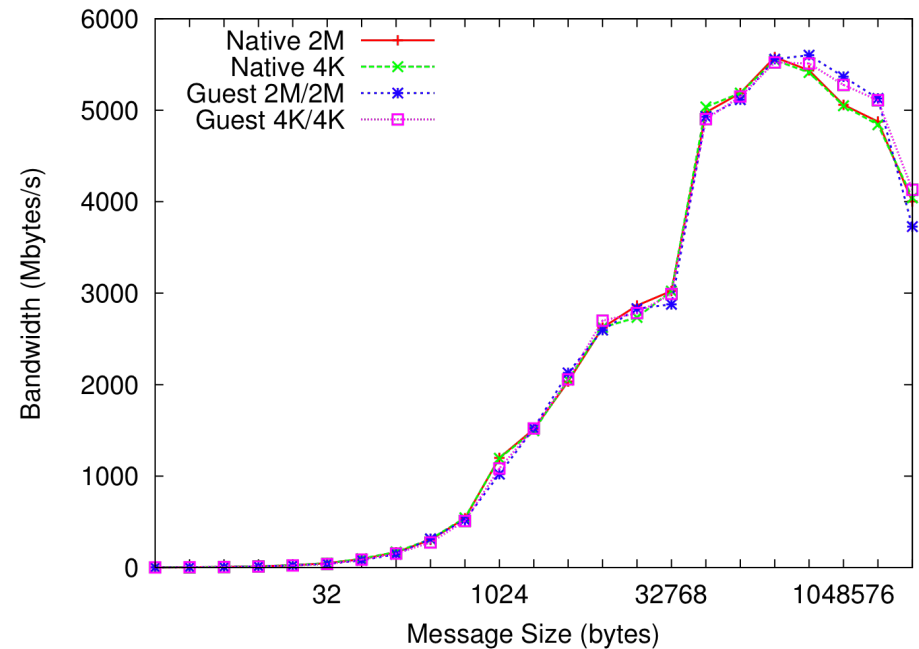
Latency in Guest More Variable

Bandwidth Essentially Identical

Latency



Bandwidth

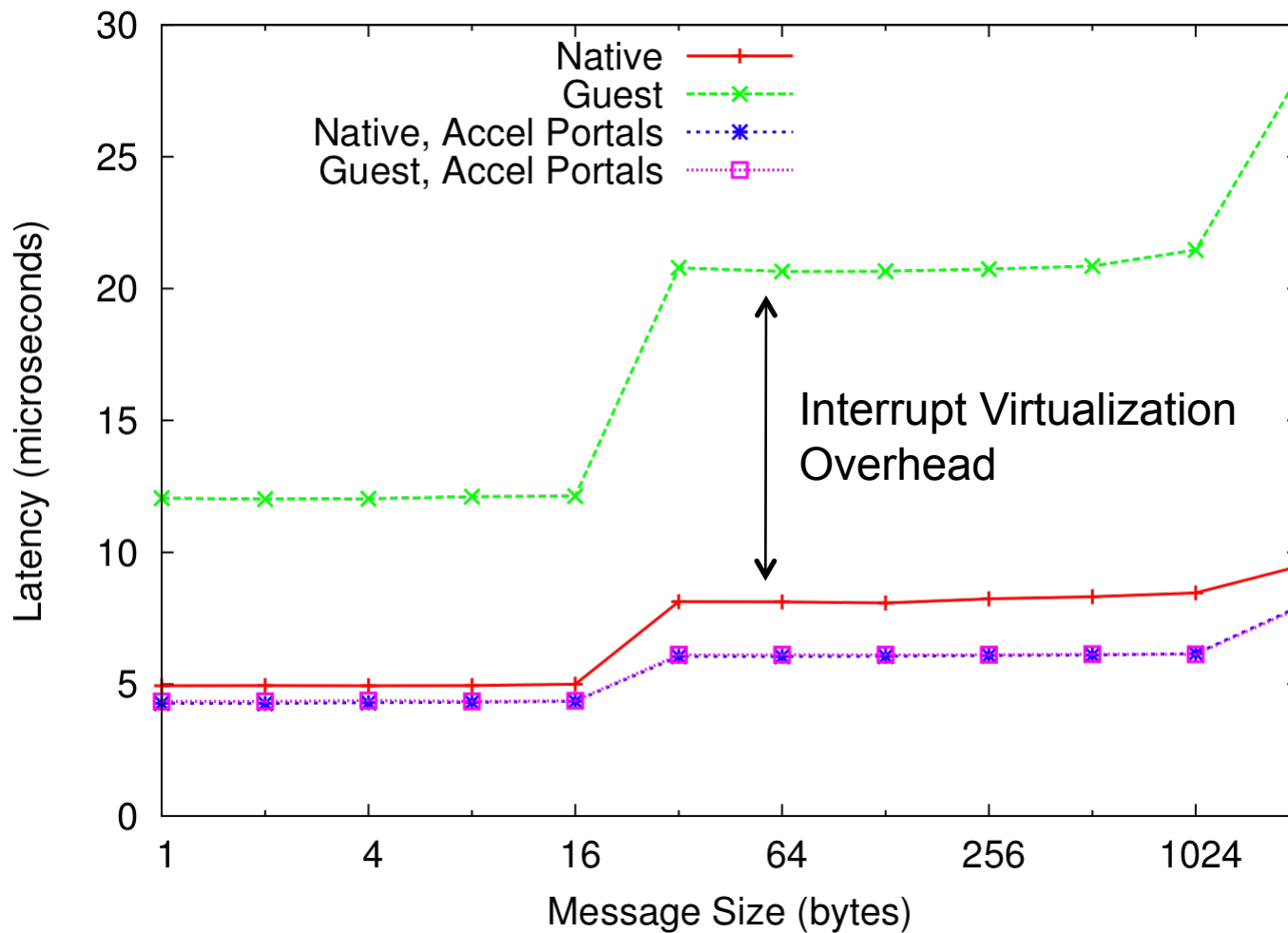


Variability possibly due to timekeeping inaccuracy in guest



VMM-Bypass MPI Latency on Cray XT4

Avoiding Interrupt Virtualization Important





Conclusions

- **Virtualization support continuously improving**
- **Significant previous HPC virtualization work**
- **Compelling use cases for high-end HPC**
 - Increase flexibility
 - Enable new capabilities
- **Results show low virtualization overhead**
 - NUMA and VCPU pinning important in all cases
 - Large paging important for random access



Acknowledgements

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 - DOE ASCR X-stack (current)
 - DOE ASC (current)
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